

Interpolation Properties of Proofs with Cuts

Anela Lolić
TU Wien

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Maehara's partition method

Definition (LK#)

Gentzen's system LK + the predicate symbol \top with 0 argument places and the additional axiom $\vdash \top$.

Definition (partition)

Let $S : \Gamma \vdash \Delta$ be a sequent and let Γ_1, Γ_2 be a permutation variant of Γ , and Δ_1, Δ_2 a permutation variant of Δ . Then $[\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ is called a partition of S .

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Let $S : \Gamma \vdash \Delta$ be a sequent and let Γ_1, Γ_2 be a permutation variant of Γ , and Δ_1, Δ_2 a permutation variant of Δ . Then $[\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ is called a partition of S .

Example

Consider the sequent $\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))$.

$$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{ ; \exists y(P(a) \rightarrow Q(y))\}]$$

Weak interpolant

Definition

Let S be a sequent and $X : [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ a partition of S . Then the formula I is called a weak interpolant of S w.r.t. X if

1. φ_1 is an LK $\#$ -proof of $\Gamma_1 \vdash \Delta_1, I$, and φ_2 is an LK $\#$ -proof of $I, \Gamma_2 \vdash \Delta_2$.
2. The predicate symbols in I are a subset of the predicate symbols occurring in $\{\Gamma_1, \Delta_1\}$ and $\{\Gamma_2, \Delta_2\}$.

A proof of the form

$$\frac{\begin{array}{c} (\varphi_1) \\ \Gamma_1 \vdash \Delta_1, I \end{array} \quad \begin{array}{c} (\varphi_2) \\ I, \Gamma_2 \vdash \Delta_2 \end{array}}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \text{ cut}$$

is called an interpolation derivation for S w.r.t. X .

Strong interpolant

weak interpolant + all its **free variables**, **constant symbols** and **function symbols** occur in $\{\Gamma_1, \Delta_1\}$ and $\{\Gamma_2, \Delta_2\}$.

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weak interpolant + all its **free variables**, **constant symbols** and **function symbols** occur in $\{\Gamma_1, \Delta_1\}$ and $\{\Gamma_2, \Delta_2\}$.

Lemma

Let I be a weak interpolant of a sequent S w.r.t. a partition X . Then there exists an strong interpolant I^ of S w.r.t. X .*

Proof.

$X : [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ partition of S , and the corresponding weak interpolation derivation $\psi =$

$$\frac{\begin{array}{c} (\varphi_1) \\ \Gamma_1 \vdash \Delta_1, I \end{array} \quad \begin{array}{c} (\varphi_2) \\ I, \Gamma_2 \vdash \Delta_2 \end{array}}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \text{ cut}$$

Select a maximal term t not occurring in both partitions, assume $t = f(t_1, \dots, t_n)$ for some function symbol f not occurring as a function symbol in both partitions. We distinguish

Proof.

$X : [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ partition of S , and the corresponding weak interpolation derivation $\psi =$

$$\frac{(\varphi_1) \quad \Gamma_1 \vdash \Delta_1, I \quad I, \Gamma_2 \vdash \Delta_2 \quad (\varphi_2)}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \text{ cut}$$

Select a maximal term t not occurring in both partitions, assume $t = f(t_1, \dots, t_n)$ for some function symbol f not occurring as a function symbol in both partitions. We distinguish

1. f occurs only in $\{\Gamma_1; \Delta_1\}$.

$$\frac{\frac{(\varphi_1) \quad \Gamma_1 \vdash \Delta_1, I}{\Gamma_1 \vdash \Delta_1, \exists x I\{t \leftarrow x\}} \exists_r \quad \frac{(\varphi_2\{t \leftarrow \alpha\}) \quad I\{t \leftarrow \alpha\}, \Gamma_2 \vdash \Delta_2}{\exists x I\{t \leftarrow x\}, \Gamma_2 \vdash \Delta_2} \exists_l}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \text{ cut}$$

Proof.

$X : [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ partition of S , and the corresponding weak interpolation derivation $\psi =$

$$\frac{(\varphi_1) \quad \Gamma_1 \vdash \Delta_1, I \quad I, \Gamma_2 \vdash \Delta_2 \quad (\varphi_2)}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \text{ cut}$$

Select a maximal term t not occurring in both partitions, assume $t = f(t_1, \dots, t_n)$ for some function symbol f not occurring as a function symbol in both partitions. We distinguish

2. f occurs only in $\{\Gamma_2; \Delta_2\}$.

$$\frac{\frac{(\varphi_1 \{t \leftarrow \alpha\}) \quad \Gamma_1 \vdash \Delta_1, I \{t \leftarrow \alpha\}}{\Gamma_1 \vdash \Delta_1, \forall x I \{t \leftarrow x\}} \forall_r \quad \frac{(\varphi_2) \quad I, \Gamma_2 \vdash \Delta_2}{\forall x I \{t \leftarrow x\}, \Gamma_2 \vdash \Delta_2} \forall_l}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \text{ cut}$$

Proof.

$X : [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ partition of S , and the corresponding weak interpolation derivation $\psi =$

$$\frac{\begin{array}{c} (\varphi_1) \\ \Gamma_1 \vdash \Delta_1, I \end{array} \quad \begin{array}{c} (\varphi_2) \\ I, \Gamma_2 \vdash \Delta_2 \end{array}}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \text{ cut}$$

Select a maximal term t not occurring in both partitions, assume $t = f(t_1, \dots, t_n)$ for some function symbol f not occurring as a function symbol in both partitions. We distinguish

3. f does not occur in the partitions. Then the interpolant can be constructed as in case 1 or in case 2, as both constructions work.

Lemma (Maehara's lemma)

Let $\Gamma \vdash \Delta$ be **LK**-provable, and $X: [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2\}]$ an arbitrary partition of $\Gamma \vdash \Delta$. Then there exists a formula I , called the interpolant of $\Gamma \vdash \Delta$ w.r.t. the partition X , s.t.

1. $\Gamma_1 \vdash \Delta_1, I$ and $I, \Gamma_2 \vdash \Delta_2$ are both **LK**_#-provable.
2. I contains only free variables and individual and predicate constants apart from \top that occur in $\{\Gamma_1; \Delta_1\}$ and $\{\Gamma_2; \Delta_2\}$.

Proof.

By induction on the number of inferences k in a cut-free proof of $\Gamma \vdash \Delta$.

- ▶ Interpolants I s.t. $\Gamma \vdash I$ and $I \vdash \Delta$ are provable.
- ▶ $k = 0$. The sequent $\Gamma \vdash \Delta$ is an axiom of the form $C \vdash C$. We look at the partition $[\{C; \}, \{; C\}]$. C fulfills all requirements for an interpolant of $C \vdash C$.
- ▶ $k > 0$: Consider the last inference in the derivation.

Proof.

Let \forall_I be the last inference:

$$\frac{F(t), \Gamma \vdash \Delta}{\forall x F(x), \Gamma \vdash \Delta} \forall_I$$

with partition $[\{\forall x F(x), \Gamma; \}, \{; \Delta\}]$.

The partition in the upper sequent is $[\{F(t), \Gamma; \}, \{; \Delta\}]$.

By IH:

$$F(t), \Gamma \vdash I(b_1, \dots, b_n) \text{ and } I(b_1, \dots, b_n) \vdash \Delta.$$

b_1, \dots, b_n the free variables and constants occurring in t .

Replace the b_{i_1}, \dots, b_{i_m} that do not occur in $F(x), \Delta$ by bound variables y_1, \dots, y_m :

$$I' = \forall y_1, \dots, \forall y_m I(b_1, \dots, y_1, \dots, y_m, \dots, b_n)$$

$$\frac{\frac{\frac{F(t), \Gamma \vdash I(b_1, \dots, b_n)}{\forall x F(x), \Gamma \vdash I(b_1, \dots, b_n)} \forall_I}{\forall x F(x), \Gamma \vdash I'} \forall_r \quad \frac{I(b_1, \dots, b_n) \vdash \Delta}{I' \vdash \Delta} \forall_I}{\forall x F(x), \Gamma \vdash \Delta} \text{cut}$$

Proof.

Let \forall_r be the last inference:

$$\frac{\Gamma \vdash \Delta, F(y)}{\Gamma \vdash \Delta, \forall x F(x)} \forall_r$$

with partition $[\{\Gamma; \}, \{; \Delta, \forall x F(x)\}]$.

The partition in the upper sequent is $[\{\Gamma; \}, \{; \Delta, F(y)\}]$.

By IH:

$$\Gamma \vdash I \text{ and } I \vdash \Delta, F(y).$$

Since y does not occur in $\Gamma, \Delta, F(x)$ it does not occur in I and we infer

$$\frac{I \vdash \Delta, F(y)}{I \vdash \Delta, \forall x F(x)} \forall_r$$

I is also an interpolant for the lower sequent $\Gamma \vdash \Delta, \forall x F(x)$.

Craig's interpolation theorem

Theorem

Let A and B be formulas s.t. $A \rightarrow B$ is **LK**-provable.

If A and B have at least one predicate constant in common, then there is a formula I , called the interpolant of A and B , s.t.

- ▶ I only contains free variables and individual and predicate constants that occur in both A and B ,
- ▶ and $A \rightarrow I$ and $I \rightarrow B$ are **LK**-provable.

If A and B have no predicate constant in common, then either $A \rightarrow$ or $\rightarrow B$ is **LK**-provable.

Example

$$\frac{\frac{\frac{P(a) \vdash P(a) \quad Q(a) \vdash Q(a)}{P(a), P(a) \rightarrow Q(a) \vdash Q(a)} \rightarrow_I}{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)} \rightarrow_r}{\frac{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))} \exists_r} \forall_I$$

$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{\exists y(P(a) \rightarrow Q(y))\}]$

Example

$$\frac{\frac{\frac{P(a) \vdash P(a) \quad Q(a) \vdash Q(a)}{P(a), P(a) \rightarrow Q(a) \vdash Q(a)} \rightarrow_I}{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)} \rightarrow_r}{\frac{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))} \exists_r} \forall_I$$

$$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{\exists y(P(a) \rightarrow Q(y))\}]$$

axioms:

$$[\{\vdash P(a)\}, \{P(a)\}] \text{ and } [\{\vdash Q(a)\}, \{Q(a)\}]$$

Example

$$\frac{\frac{\frac{P(a) \vdash P(a) \quad Q(a) \vdash Q(a)}{P(a), P(a) \rightarrow Q(a) \vdash Q(a)} \rightarrow_I}{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)} \rightarrow_r}{\frac{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))} \exists_r} \forall_I$$

$$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{\exists y(P(a) \rightarrow Q(y))\}]$$

axioms:

$$[\{\vdash P(a)\}, \{P(a)\}] \text{ and } [\{Q(a)\}, \{Q(a)\}]$$

$$\vdash P(a), \neg P(a) \text{ and } \neg P(a), P(a) \vdash$$

Example

$$\begin{array}{c}
 \frac{P(a) \vdash P(a) \quad Q(a) \vdash Q(a)}{P(a), P(a) \rightarrow Q(a) \vdash Q(a)} \rightarrow_I \\
 \frac{P(a), P(a) \rightarrow Q(a) \vdash Q(a)}{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)} \rightarrow_r \\
 \frac{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)}{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))} \exists_r \\
 \frac{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))} \forall_I
 \end{array}$$

$$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{\exists y(P(a) \rightarrow Q(y))\}]$$

axioms:

$$[\{\vdash P(a)\}, \{P(a)\}] \text{ and } [\{\vdash Q(a)\}, \{Q(a)\}]$$

$$\vdash P(a), \neg P(a) \text{ and } \neg P(a), P(a) \vdash \quad \text{and} \quad Q(a) \vdash Q(a) \text{ and } Q(a) \vdash Q(a)$$

Example

$$\begin{array}{c}
 \frac{P(a) \vdash P(a) \quad Q(a) \vdash Q(a)}{P(a), P(a) \rightarrow Q(a) \vdash Q(a)} \rightarrow_I \\
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 \frac{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)}{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))} \exists_r \\
 \frac{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))} \forall_I
 \end{array}$$

$$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{\}; \{\exists y(P(a) \rightarrow Q(y))\}]$$

axioms:

$$[\{\}; P(a)], \{P(a); \} \quad \text{and} \quad [\{Q(a); \}, \{\}; Q(a)]$$

$$\vdash P(a), \neg P(a) \quad \text{and} \quad \neg P(a), P(a) \vdash \quad \text{and} \quad Q(a) \vdash Q(a) \quad \text{and} \quad Q(a) \vdash Q(a)$$

$$\neg P(a) \vee Q(a)$$

Example

$$\begin{array}{c}
 \frac{P(a) \vdash P(a) \quad Q(a) \vdash Q(a)}{P(a), P(a) \rightarrow Q(a) \vdash Q(a)} \rightarrow_I \\
 \frac{P(a), P(a) \rightarrow Q(a) \vdash Q(a)}{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)} \rightarrow_r \\
 \frac{P(a) \rightarrow Q(a) \vdash P(a) \rightarrow Q(a)}{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))} \exists_r \\
 \frac{P(a) \rightarrow Q(a) \vdash \exists y(P(a) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))} \forall_I
 \end{array}$$

$$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{\}; \{\exists y(P(a) \rightarrow Q(y))\}]$$

axioms:

$$[\{\}; P(a)], \{P(a); \} \quad \text{and} \quad [\{Q(a); \}, \{\}; Q(a)]$$

$$\vdash P(a), \neg P(a) \quad \text{and} \quad \neg P(a), P(a) \vdash \quad \text{and} \quad Q(a) \vdash Q(a) \quad \text{and} \quad Q(a) \vdash Q(a)$$

$$\neg P(a) \vee Q(a)$$

$$\forall x(\neg P(x) \vee Q(x))$$

The case of atomic cuts

Lemma

Let φ be an LK-proof of the form

$$\frac{(\varphi_1) \quad (\varphi_2)}{\Gamma \vdash \Delta, F \quad F, \Pi \vdash \Lambda} \text{ cut}$$
$$\Gamma, \Pi \vdash \Delta, \Lambda$$

F atomic,

I interpolant of $\Gamma \vdash \Delta, F$,

J interpolant of $F, \Pi \vdash \Lambda$.

Then there exists an interpolant of the end-sequent of the form

$$I \wedge J \text{ or } I \vee J.$$

Proof.

Let $X = [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$. We distinguish:

1. F occurs only in $\{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}$. Then,

$$X_1 = [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2, F\}] \text{ of } \Gamma \vdash \Delta, F$$

$$X_2 = [\{\Pi_1; \Lambda_1\}, \{F, \Pi_2; \Lambda_2\}] \text{ of } F, \Pi \vdash \Lambda.$$

Then there are interpolation derivations $\varphi_1 =$

$$\frac{\begin{array}{c} (\chi_{1,1}) \\ \Gamma_1 \vdash \Delta_1, I \end{array} \quad \begin{array}{c} (\chi_{1,2}) \\ I, \Gamma_2 \vdash \Delta_2, F \end{array}}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2, F} \text{ cut}$$

$\varphi_2 =$

$$\frac{\begin{array}{c} (\chi_{2,1}) \\ \Pi_1 \vdash \Lambda_1, J \end{array} \quad \begin{array}{c} (\chi_{2,2}) \\ J, F, \Pi_2 \vdash \Lambda_2 \end{array}}{\Pi_1, \Pi_2, F \vdash \Lambda_1, \Lambda_2} \text{ cut}$$

Proof.

Let $X = [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$. We distinguish:

1. F occurs only in $\{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}$. Then,

$$X_1 = [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2, F\}] \text{ of } \Gamma \vdash \Delta, F$$

$$X_2 = [\{\Pi_1; \Lambda_1\}, \{F, \Pi_2; \Lambda_2\}] \text{ of } F, \Pi \vdash \Lambda.$$

Then there are interpolation derivations $\varphi_1 =$

$$\frac{\begin{array}{c} (\chi_{1,1}) \\ \Gamma_1 \vdash \Delta_1, I \end{array} \quad \begin{array}{c} (\chi_{1,2}) \\ I, \Gamma_2 \vdash \Delta_2, F \end{array}}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2, F} \text{ cut}$$

$\varphi_2 =$

$$\frac{\begin{array}{c} (\chi_{2,1}) \\ \Pi_1 \vdash \Lambda_1, J \end{array} \quad \begin{array}{c} (\chi_{2,2}) \\ J, F, \Pi_2 \vdash \Lambda_2 \end{array}}{\Pi_1, \Pi_2, F \vdash \Lambda_1, \Lambda_2} \text{ cut}$$

Proof.

Let $X = [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$. We distinguish:

1. F occurs only in $\{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}$. Then,

$$X_1 = [\{\Gamma_1; \Delta_1\}, \{\Gamma_2; \Delta_2, F\}] \text{ of } \Gamma \vdash \Delta, F$$

$$X_2 = [\{\Pi_1; \Lambda_1\}, \{F, \Pi_2; \Lambda_2\}] \text{ of } F, \Pi \vdash \Lambda.$$

$$\frac{\frac{\frac{(\chi_{1,1})}{\Gamma_1 \vdash \Delta_1, I} \quad \frac{(\chi_{2,1})}{\Pi_1 \vdash \Lambda_1, J}}{\Gamma_1, \Pi_1 \vdash \Delta_1, \Lambda_1, I \wedge J} \wedge_r \quad \frac{\frac{(\chi_{1,2})}{I, \Gamma_2 \vdash \Delta_2, F} \quad \frac{(\chi_{2,2})}{J, F, \Pi_2 \vdash \Lambda_2}}{I \wedge J, \Gamma_2, \Pi_2 \vdash \Delta_2, \Lambda_2} cut + \wedge_l}{\Gamma_1, \Pi_1, \Gamma_2, \Pi_2 \vdash \Delta_1, \Lambda_1, \Delta_2, \Lambda_2} cut$$

Proof.

Let $X = [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$. We distinguish:

2. F occurs only in $\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}$. Then,

$$X_1 = [\{\Gamma_1; \Delta_1, F\}, \{\Gamma_2; \Delta_2\}] \text{ of } \Gamma \vdash \Delta, F$$

$$X_2 = [\{F, \Pi_1; \Lambda_1\}, \{\Pi_2; \Lambda_2\}] \text{ of } F, \Pi \vdash \Lambda.$$

Then there are interpolation derivations $\varphi_1 =$

$$\frac{\begin{array}{c} (\chi_{1,1}) \\ \Gamma_1 \vdash \Delta_1, F, I \end{array} \quad \begin{array}{c} (\chi_{1,2}) \\ I, \Gamma_2 \vdash \Delta_2 \end{array}}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2, F} \text{ cut}$$

$\varphi_2 =$

$$\frac{\begin{array}{c} (\chi_{2,1}) \\ F, \Pi_1 \vdash \Lambda_1, J \end{array} \quad \begin{array}{c} (\chi_{2,2}) \\ J, \Pi_2 \vdash \Lambda_2 \end{array}}{\Pi_1, \Pi_2, F \vdash \Lambda_1, \Lambda_2} \text{ cut}$$

Proof.

Let $X = [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$. We distinguish:

2. F occurs only in $\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}$. Then,

$$X_1 = [\{\Gamma_1; \Delta_1, F\}, \{\Gamma_2; \Delta_2\}] \text{ of } \Gamma \vdash \Delta, F$$

$$X_2 = [\{F, \Pi_1; \Lambda_1\}, \{\Pi_2; \Lambda_2\}] \text{ of } F, \Pi \vdash \Lambda.$$

Then there are interpolation derivations $\varphi_1 =$

$$\frac{\begin{array}{c} (\chi_{1,1}) \\ \Gamma_1 \vdash \Delta_1, F, I \end{array} \quad \begin{array}{c} (\chi_{1,2}) \\ I, \Gamma_2 \vdash \Delta_2 \end{array}}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2, F} \text{ cut}$$

$\varphi_2 =$

$$\frac{\begin{array}{c} (\chi_{2,1}) \\ F, \Pi_1 \vdash \Lambda_1, J \end{array} \quad \begin{array}{c} (\chi_{2,2}) \\ J, \Pi_2 \vdash \Lambda_2 \end{array}}{\Pi_1, \Pi_2, F \vdash \Lambda_1, \Lambda_2} \text{ cut}$$

Proof.

Let $X = [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$. We distinguish:

2. F occurs only in $\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}$. Then,

$$X_1 = [\{\Gamma_1; \Delta_1, F\}, \{\Gamma_2; \Delta_2\}] \text{ of } \Gamma \vdash \Delta, F$$

$$X_2 = [\{F, \Pi_1; \Lambda_1\}, \{\Pi_2; \Lambda_2\}] \text{ of } F, \Pi \vdash \Lambda.$$

$$\frac{\frac{\frac{(\chi_{1,1})}{\Gamma_1 \vdash \Delta_1, F, I} \quad \frac{(\chi_{2,1})}{F, \Pi_1 \vdash \Lambda_1, J}}{\Gamma_1, \Pi_1 \vdash \Delta_1, \Lambda_1, I \vee J} \text{ cut} + \vee_r \quad \frac{\frac{(\chi_{1,2})}{I, \Gamma_2 \vdash \Delta_2} \quad \frac{(\chi_{2,2})}{J, \Pi_2 \vdash \Lambda_2}}{I \vee J, \Gamma_2, \Pi_2 \vdash \Delta_2, \Lambda_2} \vee_l}{\Gamma_1, \Pi_1, \Gamma_2, \Pi_2 \vdash \Delta_1, \Lambda_1, \Delta_2, \Lambda_2} \text{ cut}$$

Proof.

Let $X = [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$. We distinguish:

3. F does not occur in any of the partitions. Then both constructions from above work.

How about more complex cuts?

Definition

Let φ be an LK-proof of the form

$$\frac{\begin{array}{c} (\varphi_1) \\ \Gamma \vdash \Delta, F \end{array} \quad \begin{array}{c} (\varphi_2) \\ F, \Pi \vdash \Lambda \end{array}}{\Gamma, \Pi \vdash \Delta, \Lambda} \text{ cut}$$

F contains predicate symbols $P_1, \dots, P_k, P_l, \dots, P_n$.

Arbitrary partition $X: [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$.

F is **X-violating** if a subset of P_1, \dots, P_k occurs only in $\Gamma_1, \Pi_1, \Delta_1, \Lambda_1$, and a subset of P_l, \dots, P_n occurs only in $\Gamma_2, \Pi_2, \Delta_2, \Lambda_2$.

F is **X-admissible** otherwise.

Lemma

Let φ be an LK-proof of the form

$$\frac{\begin{array}{c} (\varphi_1) \\ \Gamma \vdash \Delta, F \end{array} \quad \begin{array}{c} (\varphi_2) \\ F, \Pi \vdash \Lambda \end{array}}{\Gamma, \Pi \vdash \Delta, \Lambda} \text{ cut}$$

$X: [\{\Gamma_1, \Pi_1; \Delta_1, \Lambda_1\}, \{\Gamma_2, \Pi_2; \Delta_2, \Lambda_2\}]$ a partition of the end-sequent s.t.
 F is X -admissible.

I is an interpolant of $\Gamma \vdash \Delta, F$, J is an interpolant of $F, \Pi \vdash \Lambda$.

Then there exists an interpolant of S w.r.t. X of the form

$$I \wedge J \quad \text{or} \quad I \vee J.$$

Example

$$\begin{array}{c}
 \frac{P(u) \vdash P(u) \quad Q(u) \vdash Q(u)}{P(u), P(u) \rightarrow Q(u) \vdash Q(u)} \rightarrow: l \\
 \frac{P(u), P(u) \rightarrow Q(u) \vdash Q(u)}{P(u) \rightarrow Q(u) \vdash P(u) \rightarrow Q(u)} \rightarrow: r \\
 \frac{P(u) \rightarrow Q(u) \vdash P(u) \rightarrow Q(u)}{P(u) \rightarrow Q(u) \vdash \exists y(P(u) \rightarrow Q(y))} \exists: r \\
 \frac{P(u) \rightarrow Q(u) \vdash \exists y(P(u) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(u) \rightarrow Q(y))} \forall: l \\
 \frac{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(u) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \forall x \exists y(P(x) \rightarrow Q(y))} \forall: r \\
 \frac{\forall x(P(x) \rightarrow Q(x)) \vdash \forall x \exists y(P(x) \rightarrow Q(y))}{\forall x(P(x) \rightarrow Q(x)) \vdash \exists y(P(a) \rightarrow Q(y))} \text{cut}
 \end{array}$$

$$X: [\{\forall x(P(x) \rightarrow Q(x)); \}, \{ ; \exists y(P(a) \rightarrow Q(y))\}]$$

Cut-formula $\forall x \exists y(P(x) \rightarrow Q(y))$ is X -admissible:

$$[\{\forall x(P(x) \rightarrow Q(x)); \forall x \exists y(P(x) \rightarrow Q(y))\}, \{ ; \}] \text{ and}$$

$$[\{\forall x \exists y(P(x) \rightarrow Q(y)); \}, \{ ; \exists y(P(a) \rightarrow Q(y))\}]$$

$$I = (\perp \vee \perp) \vee (\neg P(a) \vee Q(v)) = P(a) \rightarrow Q(v)$$

\Downarrow

$$\forall x \exists y(P(x) \rightarrow Q(y))$$